## **The PC Parallel Ports**

# **Chapter 21**

The original IBM PC design provided support for three parallel printer ports that IBM designated LPT1:, LPT2:, and LPT3:<sup>1</sup>. IBM probably envisioned machines that could support a standard dot matrix printer, a daisy wheel printer, and maybe some other auxiliary type of printer for different purposes, all on the same machine (laser printers were still a few years in the future at that time). Surely IBM did not anticipate the general use that parallel ports have received or they would probably have designed them differently. Today, the PC's parallel port controls keyboards, disk drives, tape drives, SCSI adapters, ethernet (and other network) adapters, joystick adapters, auxiliary keypad devices, other miscellaneous devices, and, oh yes, printers. This chapter will not attempt to describe how to use the parallel port for all these various purposes – this book is long enough already. However, a thorough discussion of how the parallel interface controls a printer and one other application of the parallel port (cross machine communication) should provide you with enough ideas to implement the next great parallel device.

### 21.1 Basic Parallel Port Information

There are two basic data transmission methods modern computes employ: parallel data transmission and serial data transmission. In a serial data transmission scheme (see "The PC Serial Ports" on page 1223) one device sends data to another a single bit at a time across one wire. In a parallel transmission scheme, one device sends data to another several bits at a time (in parallel) on several different wires. For example, the PC's parallel port provides eight data lines compared to the serial port's single data line. Therefore, it would seem that the parallel port would be able to transmit data eight times as fast since there are eight times as many wires in the cable. Likewise, it would seem that a serial cable, for the same price as a parallel cable, would be able to go eight times as far since there are fewer wires in the cable. And these are the common trade-offs typically given for parallel vs. serial communication methods: speed vs. cost.

In practice, parallel communications is not eight times faster than serial communications, nor do parallel cables cost eight times as much. In generally, those who design serial cables (.e.g, ethernet cables) use higher materials and shielding. This raises the cost of the cable, but allows devices to transmit data, still a bit at a time, much faster. Furthermore, the better cable design allows greater distances between devices. Parallel cables, on the other hand, are generally quite inexpensive and designed for very short connections (generally no more than about six to ten feet). The real world problems of electrical noise and cross-talk create problems when using long parallel cables and limit how fast the system can transmit data. In fact the original Centronics printer port specification called for no more than 1,000 characters/second data transmission rate, so many printers were designed to handle data at this transmission rate. Most parallel ports can easily outperform this value; however, the limiting factor is still the cable, not any intrinsic limitation in a modern computer.

Although a parallel communication system could use any number of wires to transmit data, most parallel systems use eight data lines to transmit a byte at a time. There are a few notable exceptions. For example, the SCSI interface is a parallel interface, yet newer versions of the SCSI standard allow eight, sixteen, and even thirty-two bit data transfers. In this chapter we will concentrate on byte-sized transfers since the parallel port on the PC provides for eight-bit data.

A typical parallel communication system can be one way (or *unidirectional*) or two way (*bidirectional*). The PC's parallel port generally supports unidirectional communications (from the PC to the printer), so we will consider this simpler case first.

In a unidirectional parallel communication system there are two distinguished sites: the transmitting site and the receiving site. The transmitting site places its data on the data lines and informs the receiving site that data is available; the receiving site then reads the data lines and informs the transmitting site that it

<sup>1.</sup> In theory, the BIOS allows for a fourth parallel printer port, LPT4:, but few (if any) adapter cards have ever been built that claim to work as LPT4:.

has taken the data. Note how the two sites synchronize their access to the data lines – the receiving site does not read the data lines until the transmitting site tells it to, the transmitting site does not place a new value on the data lines until the receiving site removes the data and tells the transmitting site that it has the data. *Handshaking* is the term that describes how these two sites coordinate the data transfer.

To properly implement handshaking requires two additional lines. The *strobe* (or data strobe) line is what the transmitting site uses to tell the receiving site that data is available. The *acknowledge* line is what the receiving site uses to tell the transmitting site that it has taken the data and is ready for more. The PC's parallel port actually provides a third handshaking line, *busy*, that the receiving site can use to tell the transmitting site should not attempt to send data. A typical data transmission session looks something like the following:

Transmitting site:

- 1) The transmitting site checks the busy line to see if the receiving is busy. If the busy line is active, the transmitter waits in a loop until the busy line becomes inactive.
- 2) The transmitting site places its data on the data lines.
- 3) The transmitting site activates the strobe line.
- 4) The transmitting site waits in a loop for the acknowledge line to become active.
- 5) The transmitting site sets the strobe inactive.
- 6) The transmitting site waits in a loop for the acknowledge line to become inactive.
- 7) The transmitting site repeats steps one through six for each byte it must transmit.

Receiving site:

- 1) The receiving site sets the busy line inactive (assuming it is ready to accept data).
- 2) The receiving site waits in a loop until the strobe line becomes active.
- 3) The receiving site reads the data from the data lines (and processes the data, if necessary).
- 4) The receiving site activates the acknowledge line.
- 5) The receiving site waits in a loop until the strobe line goes inactive.
- 6) The receiving site sets the acknowledge line inactive.
- 7) The receiving site repeats steps one through six for each additional byte it must receive.

By carefully following these steps, the receiving and transmitting sites carefully coordinate their actions so the transmitting site doesn't attempt to put several bytes on the data lines before the receiving site consumes them and the receiving site doesn't attempt to read data that the transmitting site has not sent.

Bidirectional data transmission is often nothing more than two unidirectional data transfers with the roles of the transmitting and receiving sites reversed for the second communication channel. Some PC parallel ports (particularly on PS/2 systems and many notebooks) provide a bidirectional parallel port. Bidirectional data transmission on such hardware is slightly more complex than on systems that implement bidirectional communication with two unidirectional ports. Bidirectional communication on a bidirectional ports an extra set of control lines so the two sites can determine who is writing to the common data lines at any one time.

## 21.2 The Parallel Port Hardware

The standard unidirectional parallel port on the PC provides more than the 11 lines described in the previous section (eight data, three handshake). The PC's parallel port provides the following signals:

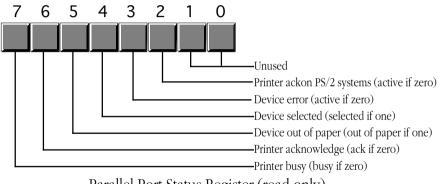
Pin Number on	I/O	Active	Signal
Connector	Direction	Polarity	Description
1	output	0	Strobe (data available signal).
2-9	output	-	Data lines (bit 0 is pin 2, bit 7 is pin 9).
10	input	0	Acknowledge line (active when remote system has taken data).
11	input	0	Busy line (when active, remote system is busy and cannot accept data).
12	input	1	Out of paper (when active, printer is out of paper).
13	input	1	Select. When active, the printer is selected.
14	output	0	Autofeed. When active, the printer automatically inserts a line feed after every carriage return it receives.
15	input	0	Error. When active, there is a printer error.
16	output	0	Init. When held active for at least 50 $\mu$ sec, this signal causes the printer to initialize itself.
17	output	0	Select input. This signal, when inactive, forces the printer off-line
18-25	-	-	Signal ground.

**Table 79: Parallel Port Signals** 

Note that the parallel port provides 12 output lines (eight data lines, strobe, autofeed, init, and select input) and five input lines (acknowledge, busy, out of paper, select, and error). Even though the port is unidirectional, there is a good mixture of input and output lines available on the port. Many devices (like disk and tape drives) that require bidirectional data transfer use these extra lines to perform bidirectional data transfer.

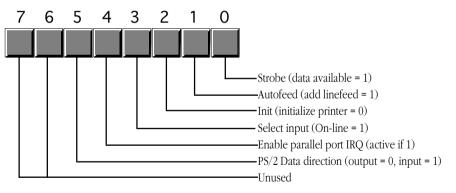
On bidirectional parallel ports (found on PS/2 and laptop systems), the strobe and data lines are both input and output lines. There is a bit in a control register associated with the parallel port that selects the transfer direction at any one given instant (you cannot transfer data in both direction simultaneously).

There are three I/O addresses associated with a typical PC compatible parallel port. These addresses belong to the *data register, the status register,* and *the control register*. The data register is an eight-bit read/write port. Reading the data register (in a unidirectional mode) returns the value last written to the data register. The control and status registers provide the interface to the other I/O lines. The organization of these ports is as follows:



Parallel Port Status Register (read only)

Bit two (printer acknowledge) is available only on PS/2 and other systems that support a bidirectional printer port. Other systems do not use this bit.



Parallel Port Control Register

The parallel port control register is an output register. Reading this location returns the last value written to the control register *except for bit five* that is write only. Bit five, the data direction bit, is available only on PS/2 and other systems that support a bidirectional parallel port. If you write a zero to this bit, the strobe and data lines are output bits, just like on the unidirectional parallel port. If you write a one to this bit, then the data and strobe lines are inputs. Note that in the input mode (bit 5 = 1), bit zero of the control register is actually an input. Note: writing a one to bit four of the control register enables the printer IRQ (IRQ 7). However, this feature does not work on all systems so very few programs attempt to use interrupts with the parallel port. When active, the parallel port will generate an int 0Fh whenever the printer acknowledges a data transmission.

Since the PC supports up to three separate parallel ports, there could be as many as three sets of these parallel port registers in the system at any one time. There are three *parallel port base addresses* associated with the three possible parallel ports: 3BCh, 378h, and 278h. We will refer to these as the base addresses for LPT1:, LPT2:, and LPT3:, respectively. The parallel port data register is always located at the base address for a parallel port, the status register appears at the base address plus one, and the control register appears at the base address plus one, and the control s3BCh, the status register is at I/O address 3BCh, and the control register is at I/O address 3BCh.

There is one minor glitch. The I/O addresses for LPT1:, LPT2:, and LPT3: given above are the *physical addresses* for the parallel ports. The BIOS provides *logical addresses* for these parallel ports as well. This lets users remap their printers (since most software only writes to LPT1:). To accomplish this, the BIOS reserves eight bytes in the BIOS variable space (40:8, 40:0A, 40:0C, and 40:0E). Location 40:8 contains the base address for logical LPT1:, location 40:0A contains the base address for logical LPT2:, etc., it generally accesses the parallel port whose base address appears in one of these locations.

## 21.3 Controlling a Printer Through the Parallel Port

Although there are many devices that connect to the PC's parallel port, printers still make up the vast number of such connections. Therefore, describing how to control a printer from the PC's parallel port is probably the best first example to present. As with the keyboard, your software can operate at three different levels: it can print data using DOS, using BIOS, or by writing directly to the parallel port hardware. As with the keyboard interface, using DOS or BIOS is the best approach if you want to maintain compatibility with other devices that plug into the parallel port<sup>2</sup>. Of course, if you are controlling some other type of

<sup>2.</sup> Many devices connect to the parallel port with a pass-through plug allowing you to use that device and still use the parallel port for your printer. However, if you talk directly to the parallel port with your software, it may conflict with that device's operation.

device, going directly to the hardware is your only choice. However, the BIOS provides good printer support, so going directly to the hardware is rarely necessary if you simply want to send data to the printer.

#### 21.3.1 Printing via DOS

MS-DOS provides two calls you can use to send data to the printer. DOS function 05h writes the character in the **d1** register directly to the printer. Function 40h, with a file handle of 04h, also sends data to the printer. Since the chapter on DOS and BIOS fully describes these functions, we will not discuss them any further here. For more information, see "MS-DOS, PC-BIOS, and File I/O" on page 699.

#### 21.3.2 Printing via BIOS

Although DOS provides a reasonable set of functions to send characters to the printer, it does not provide functions to let you initialize the printer or obtain the current printer status. Furthermore, DOS only prints to LPT1:. The PC's int 17h BIOS routine provides three functions, print, initialize, and status. You can apply these functions to any supported parallel port on the system. The print function is roughly equivalent to DOS' print character function. The initialize function initializes the printer using system dependent timing information. The printer status returns the information from the printer status port along with time-out information. For more information on these routines, see "MS-DOS, PC-BIOS, and File I/O" on page 699.

#### 21.3.3 An INT 17h Interrupt Service Routine

Perhaps the best way to see how the BIOS functions operate is to write a replacement int 17h ISR for a printer. This section explains the handshaking protocol and variables the printer driver uses. It also describes the operation and return results associated with each machine.

There are eight variables in the BIOS variable space (segment 40h) the printer driver uses. The following table describes each of these variables:

Address	Description
40:08	Base address of LPT1: device.
40:0A	Base address of LPT2: device.
40:0C	Base address of LPT3: device.
40:0E	Base address of LPT4: device.
40:78	LPT1: time-out value. The printer port driver software should return an error if the printer device does not respond in a reasonable amount of time. This variable (if non-zero) determines how many loops of 65,536 iterations each a driver will wait for a printer acknowledge. If zero, the driver will wait forever.
40:79	LPT2: time-out value. See description above.
40:7A	LPT3: time-out value. See description above.
40:7B	LPT4: time-out value. See description above.

#### **Table 80: BIOS Parallel Port Variables**

You will notice a slight deviation in the handshake protocol in the following code. This printer driver does not wait for an acknowledge from the printer *after* sending a character. Instead, it checks to see if

; INT17.ASM

the printer has sent an acknowledge to the previous character *before* sending a character. This saves a small amount of time because the program printer then characters can continue to operating in parallel with the receipt of the acknowledge from the printer. You will also notice that this particular driver does not monitor the busy lines. Almost every printer in existence leaves this line inactive (not busy), so there is no need to check it. If you encounter a printer than does manipulate the busy line, the modification to this code is trivial. The following code implements the int 17h service:

; A short passive TSR that replaces the BIOS' int 17h handler. ; This routine demonstrates the function of each of the int 17h functions that a standard BIOS would provide. : : Note that this code does not patch into int 2Fh (multiplex interrupt) ; nor can you remove this code from memory except by rebooting. ; If you want to be able to do these two things (as well as check for ; a previous installation), see the chapter on resident programs. Such ; code was omitted from this program because of length constraints. ; cseq and EndResident must occur before the standard library segments! para public `code' seament csea csea ends ; Marker segment, to find the end of the resident section. EndResident segment para public 'Resident' EndResident ends .xlist include stdlib.a includelib stdlib.lib .list byp <byte ptr> equ cseq segment para public 'code' assume cs:cseq, ds:cseq OldInt17 dword ? ; BIOS variables: PrtrBase equ 8 PrtrTimeOut equ 78h ; This code handles the INT 17H operation. INT 17H is the BIOS routine ; to send data to the printer and report on the printer's status. There ; are three different calls to this routine, depending on the contents ; of the AH register. The DX register contains the printer port number. ; DX=0 -- Use LPT1: ; DX=1 -- Use LPT2: ; DX=2 -- Use LPT3: ; DX=3 -- Use LPT4: ; ; AH=0 --Print the character in AL to the printer. Printer status is returned in AH. If bit #0 = 1 then a timeout error occurred. ; ; ; AH=1 --Initialize printer. Status is returned in AH. ; AH=2 --Return printer status in AH. The status bits returned in AH are as follows: ;

; Bit Function Non-error values ; ----\_\_\_\_\_ ; 0 1=time out error 0 ; 1 unused x ; 2 unused x 3 1=I/O error 0 : 1=selected, 0=deselected. 4 1 ; 5 1=out of paper 0 ; 6 1=acknowledge ; x 7 1=not busv x ; ; Note that the hardware returns bit 3 with zero if an error has occurred, ; with one if there is no error. The software normally inverts this bit ; before returning it to the caller. ; Printer port hardware locations: ; There are three ports used by the printer hardware: ; PrtrPortAdrs ---Output port where data is sent to printer (8 bits). ; PrtrPortAdrs+1 ---Input port where printer status can be read (8 bits). : PrtrPortAdrs+2 ---Output port where control information is sent to the printer. ; Data output port- 8-bit data is transmitted to the printer via this port. ; Input status port: bit 0: unused bit 1: ; unused. bit 2: unused. ; ; bit 3: -Error, normally this bit means that the ; printer has encountered an error. However, ; with the P101 installed this is a data ; return line for the keyboard scan. ; ; bit 4: +SLCT, normally this bit is used to determine ; if the printer is selected or not. With the ; P101 installed this is a data return ; line for the keyboard scan. ; ; bit 5: +PE, a 1 in this bit location means that the ; printer has detected the end of paper. On ; many printer ports, this bit has been found ; ; to be inoperative. ; bit 6: -ACK, A zero in this bit position means that ; the printer has accepted the last character ; and is ready to accept another. This bit ; is not normally used by the BIOS as bit 7 ; also provides this function (and more). ; ; bit 7: -Busy, When this signal is active (0) the ; printer is busy and cannot accept data. ; When this bit is set to one, the printer ; can accept another character. ; ; ; ; Output control port: Bit 0: +Strobe, A 0.5 us (minimum) active high pulse ; on this bit clocks the data latched into the ; printer data output port to the printer. ; ; +Auto FD XT - A 1 stored at this bit causes Bit 1: ; the printer to line feed after a line is ; printed. On some printer interfaces (e.g., ; ; the Hercules Graphics Card) this bit is inoperative. ; ; Bit 2: -INIT, a zero on this bit (for a minimum of ; 50 us) will cause the printer to (re)init-;

;		ialize itself.		
;				
; ; ;	Bit 3:	+SLCT IN, a one in this bit selects the printer. A zero will cause the printer to go off-line.		
; ; ;	Bit 4:	+IRQ ENABLE, a one in this bit position allows an interrupt to occur when -ACK changes from one to zero.		
; ; ;	Bit 5:	Direction control 1=input.	on BI-DIR port. 0=output,	
;	Bit 6:	reserved, must be		
;	Bit 7:	reserved, must be	zero.	
MyInt17	proc assume	far ds:nothing		
	push	ds		
	push push	bx cx		
	push	dx		
	mov	bx, 40h	;Point DS at BIOS vars.	
	mov	ds, bx		
	cmp	dx, 3	;Must be LPT1LPT4.	
	ja	InvalidPrtr		
	cmp	ah, 0	;Branch to the appropriate code for	
	jz cmp	PrtChar ah, 2	; the printer function	
	jb	PrtrInit		
	je	PrtrStatus		
InvalidPrtr:	jmp	ISR17Done		
			it line for at least 50 us. ond 50 usec even on the fastest	
PrtrInit:	mov	bx, dx	;Get printer port value.	
iiciinic.	shl	bx, 1	;Convert to byte index.	
	mov	dx, PrtrBase[bx]	;Get printer base address.	
	test		;Does this printer exist?	
	je add	InvalidPrtr dx, 2	;Quit if no such printer. ;Point dx at control reg.	
	in	al, dx	;Read current status.	
	and	al, 11011011b	;Clear INIT/BIDIR bits.	
	out	dx, al	;Reset printer.	
PIDelay:	mov loop	cx, O PIDelay	;This will produce at least ; a 50 usec delay.	
1120101.	or	al, 100b	;Stop resetting printer.	
	out	dx, al		
	jmp	ISR17Done		
	jmp	ISR17Done		
	current pri		code reads the printer status the calling code.	
	current pri prmats the b mov	nter status. This its for return to bx, dx	the calling code. ;Get printer port value.	
; port and fo	current pri prmats the b mov shl	nter status. This its for return to bx, dx bx, 1	the calling code. ;Get printer port value. ;Convert to byte index.	
; port and fo	current pri prmats the b mov shl mov	nter status. This its for return to bx, dx bx, 1 dx, PrtrBase[bx]	the calling code. ;Get printer port value. ;Convert to byte index. ;Base address of printer port.	
; port and fo	current pri prmats the b mov shl	nter status. This its for return to bx, dx bx, 1	the calling code. ;Get printer port value. ;Convert to byte index.	
; port and fo	current pri prmats the b mov shl mov mov	nter status. This its for return to bx, dx bx, 1 dx, PrtrBase[bx] al, 00101001b	the calling code. ;Get printer port value. ;Convert to byte index. ;Base address of printer port. ;Dflt: every possible error.	
; port and fo	current pri prmats the b mov shl mov test je inc	nter status. This its for return to bx, dx bx, 1 dx, PrtrBase[bx] al, 00101001b dx, dx InvalidPrtr dx	<pre>the calling code. ;Get printer port value. ;Convert to byte index. ;Base address of printer port. ;Dflt: every possible error. ;Does this printer exist? ;Quit if no such printer. ;Point at status port.</pre>	
; port and fo	current pri prmats the b mov shl mov test je inc in	nter status. This its for return to bx, dx bx, 1 dx, PrtrBase[bx] al, 00101001b dx, dx InvalidPrtr dx al, dx	<pre>the calling code. ;Get printer port value. ;Convert to byte index. ;Base address of printer port. ;Dflt: every possible error. ;Does this printer exist? ;Quit if no such printer. ;Point at status port. ;Read status port.</pre>	
; port and fo	current pri prmats the b mov shl mov test je inc	nter status. This its for return to bx, dx bx, 1 dx, PrtrBase[bx] al, 00101001b dx, dx InvalidPrtr dx	<pre>the calling code. ;Get printer port value. ;Convert to byte index. ;Base address of printer port. ;Dflt: every possible error. ;Does this printer exist? ;Quit if no such printer. ;Point at status port.</pre>	

; Print the character in the accumulator!

PrtChar:	mov	bx, dx
	mov	<pre>cl, PrtrTimeOut[bx] ;Get time out value.</pre>
	shl	bx, 1 ;Convert to byte index.
	mov	dx, PrtrBase[bx] ;Get Printer port address
	or	dx, dx ;Non-nil pointer?
	ίz	NoPrtr2 ; Branch if a nil ptr

; The following code checks to see if an acknowlege was received from ; the printer. If this code waits too long, a time-out error is returned. ; Acknowlege is supplied in bit #7 of the printer status port (which is ; the next address after the printer data port).

	push	ax	
	inc	dx	;Point at status port
	mov	bl, cl	;Put timeout value in bl
	mov	bh, cl	; and bh.
WaitLp1:	xor	CX, CX	; Init count to 65536.
WaitLp2:	in	al, dx	;Read status port
	mov	ah, al	;Save status for now.
	test	al, 80h	;Printer acknowledge?
	jnz	GotAck	;Branch if acknowledge.
	loop	WaitLp2	;Repeat 65536 times.
	dec	bl	;Decrement time out value.
	jnz	WaitLp1	;Repeat 65536*TimeOut times.

; See if the user has selected no timeout:

cmp	bh, 0
je	WaitLpl

; TIMEOUT ERROR HAS OCCURRED!

;

Р

; A timeout - I/O error is returned to the system at this point.

; Either we fall through to this point from above (time out error) or ; the referenced printer port doesn't exist. In any case, return an error.

; the referenced printer port doesn't exist. In any case, return an error.

NoPrtr2:	or	ah, 9	;Set timeout-I/O error flags
	and	ah, 0F9h	;Turn off unused flags.
	xor	ah, 40h	;Flip busy bit.

; Okay, restore registers and return to caller.

pop	CX	;Remove old ax.
mov	al, cl	;Restore old al.
jmp	ISR17Done	

; If the printer port exists and we've received an acknowlege, then it's ; okay to transmit data to the printer. That job is handled down here.

GotAck:	mov	cx, 16	;Short delay if crazy prtr
GALp:	loop	GALp	; needs hold time after ack.
	pop	ax	;Get char to output and
	push	ax	; save again.
	dec	dx	;Point DX at printer port.
	pushf		;Turn off interrupts for now.
	cli		
	out	dx, al	;Output data to the printer.

; The following short delay gives the data time to travel through the ; parallel lines. This makes sure the data arrives at the printer before ; the strobe (the times can vary depending upon the capacitance of the ; parallel cable's lines).

	mov	cx, 16	;Give data time to settle
DataSettleLp:	loop	DataSettleLp	; before sending strobe.

; Now that the data has been latched on the printer data output port, a ; strobe must be sent to the printer. The strobe line is connected to

; bit zero of the control port. Also note that this clears bit 5 of the ; control port. This ensures that the port continues to operate as an ; output port if it is a bidirectional device. This code also clears bits ; six and seven which IBM claims should be left zero.

, SIX and Sev	en whitch it	in craims shourd be	Tert Zero.
	inc inc and out	dx dx al, dx al, 01eh dx, al	;Point DX at the printer ; control output port. ;Get current control bits. ;Force strobe line to zero and ; make sure it's an output port.
Delay0:	mov loop	cx, 16 Delay0	;Short delay to allow data ; to become good.
	or out	al, 1 dx, al	;Send out the (+) strobe. ;Output (+) strobe to bit 0
StrobeDelay:	mov loop	cx, 16 StrobeDelay	;Short delay to lengthen strobe
	and out popf	al, OFEh dx, al	;Clear the strobe bit. ;Output to control port. ;Restore interrupts.
	pop mov	dx al, dl	;Get old AX value ;Restore old AL value
ISR17Done:	pop pop pop pop iret	dx cx bx ds	
MyInt17	endp		
Main	proc		
	mov mov	ax, cseg ds, ax	
	print byte byte	"INT 17h Replacem" "Installing",	
; statements	above have tore the ol	interrupt vector. made cseg the curr d INT 17 value dir	ent data segment,
	cli mov	ax, 0	;Turn off interrupts!
	mov mov mov mov mov mov sti	es, ax ax, es:[17h*4] word ptr OldInt17 ax, es:[17h*4 + 2 word ptr OldInt17 es:[17h*4], offse es:[17h*4+2], cs	2] 7+2, ax
; We're hooke ; stay reside		only thing that rem	ains is to terminate and
	print byte	"Installed.",cr,l	.f,0
	mov int	ah, 62h 21h	;Get this program's PSP ; value.
	mov sub	dx, EndResident;C dx, bx	Compute size of program.
	mov	ax, 3100h	;DOS TSR command.

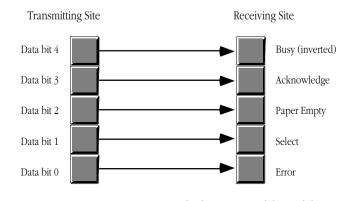
Main cseg	int endp ends	21h
sseg stk sseg	segment byte ends	para stack 'stack' 1024 dup ("stack")
zzzzzzseg LastBytes zzzzzseg	segment byte ends	para public `zzzzzz' 16 dup (?)
2	end	Main

#### 21.4 Inter-Computer Communications on the Parallel Port

Although printing is, by far, the most popular use for the parallel port on a PC, many devices use the parallel port for other purposes, as mentioned earlier. It would not be fitting to close this chapter without at least one example of a non-printer application for the parallel port. This section will describe how to get two computers to transmit files from one to the other across the parallel port.

The Laplink<sup>™</sup> program from Travelling Software is a good example of a commercial product that can transfer data across the PC's parallel port; although the following software is not as robust or feature laden as Laplink, it does demonstrate the basic principles behind such software.

Note that you cannot connect two computer's parallel ports with a simple cable that has DB25 connectors at each end. In fact, doing so could damage the computers' parallel ports because you'd be connecting digital outputs to digital outputs (a real no-no). However, you purchase "Laplink compatible" cables (or buy *real* Laplink cables for that matter) the provide proper connections between the parallel ports of two computers. As you may recall from the section on the parallel port hardware, the unidirectional parallel port provides five input signals. A Laplink cable routes four of the data lines to four of these input lines in both directions. The connections on a Laplink compatible cable are as follows:



Connections on a Laplink Compatible Cable

Data written on bits zero through three of the data register at the transmitting site appear, unchanged, on bits three through six of the status port on the receiving site. Bit four of the transmitting site appears, inverted, at bit seven of the receiving site. Note that Laplink compatible cables are bidirectional. That is, you can transmit data from either site to the other using the connections above. However, since there are only five input bits on the parallel port, you must transfer the data four bits at a time (we need one bit for the data strobe). Since the receiving site needs to acknowledge data transmissions, we cannot simultaneously transmit data in both directions. We must use one of the output lines at the site receiving data to acknowledge the incoming data.

Since the two sites cooperating in a data transfer across the parallel cable must take turns transmitting and receiving data, we must develop a *protocol* so each participant in the data transfer knows when it is okay to transmit and receive. Our protocol will be very simple – a site is either a transmitter or a receiver, the roles will never switch. Designing a more complex protocol is not difficult, but this simple protocol will suffice for the example you are about to see. Later in this section we will discuss ways to develop a protocol that allows two-way transmissions.

The following example programs will transmit and receive a single file across the parallel port. To use this software, you run the *transmit* program on the transmitting site and the *receive* program on the receiving site. The transmission program fetches a file name from the DOS command line and opens that file for reading (generating an error, and quitting, if the file does not exist). Assuming the file exists, the transmit program then queries the receiving site to see if it is available. The transmitter checks for the presence of the receiving site by alternately writing zeros and ones to all output bits then reading its input bits. The receiving site will invert these values and write them back when it comes on-line. Note that the order of execution (transmitter first or receiver first) does not matter. The two programs will attempt to hand-shake until the other comes on line. When both sites cycle through the inverting values three times, they write the value 05h to their output ports to tell the other site they are ready to proceed. A time-out function aborts either program if the other site does not respond in a reasonable amount of time.

Once the two sites are synchronized, the transmitting site determines the size of the file and then transmits the file name and size to the receiving site. The receiving site then begins waiting for the receipt of data.

The transmitting site sends the data 512 bytes at a time to the receiving site. After the transmission of 512 bytes, the receiving site delays sending an acknowledgment and writes the 512 bytes of data to the disk. Then the receiving site sends the acknowledge and the transmitting site begins sending the next 512 bytes. This process repeats until the receiving site has accepted all the bytes from the file.

Here is the code for the transmitter:

```
TRANSMIT.ASM
;
 This program is the transmitter portion of the programs that transmit files
;
; across a Laplink compatible parallel cable.
 This program assumes that the user want to use LPT1: for transmission.
:
 Adjust the equates, or read the port from the command line if this
;
; is inappropriate.
              .286
              .xlist
              include
                        stdlib.a
              includelib stdlib.lib
              .list
dsea
              segment
                        para public 'data'
                         4000
TimeOutConst equ
                                          ;About 1 min on 66Mhz 486.
                        10
                                          ;Offset to LPT1: adrs.
PrtrBase
             equ
MvPortAdrs
             word
                         ?
                                         ;Holds printer port address.
FileHandle word
                                          ;Handle for output file.
FileBuffer
             byte
                        512 dup (?) ;Buffer for incoming data.
FileSize dword ?
                                       ;Size of incoming file.
                         ?
FileNamePtr dword
                                          ;Holds ptr to filename
dseg
              ends
cseq
              segment
                        para public 'code'
                        cs:cseg, ds:dseg
              assume
; TestAbort- Check to see if the user has pressed ctrl-C and wants to
              abort this program. This routine calls BIOS to see if the
```

; ; ;		tion AH=8, read	If so, it calls DOS to read the a key w/o echo and with ctrl-C
TestAbort	proc push push mov int je mov int	near ax cx dx ah, 1 16h NoKeyPress ah, 8 21h	;See if keypress. ;Return if no keypress. ;Read char, chk for ctrl-C. ;DOS aborts if ctrl-C.
NoKeyPress:	pop pop pop ret	dx cx ax	
TestAbort	endp		
; SendByte- ;	Transmit t at a time.	-	to the receiving site four bits
SendByte	proc	near	
	push	cx	
	push mov	dx ah, al	;Save byte to xmit.
	mov	dx, MyPortAdrs	Base address of LPT1: port.
; first, just ; in the busy	bit of the	receiver.	to bit #4. This reads as a one
	mov out	al, O dx, al	;Data not ready yet.
; to bit #4 of; ; one in our l ; until the re ; one to bit	f its data busy bit (b eceiver tel #4 (which w y so often	register while it 7 of the sta ls us its ready e read as a zer	The receiver will write a zero it is busy. This comes out as a tus register). This loop waits to receive data by writing a o). Note that we check for a he user wants to abort the
	inc	dx	;Point at status register.
W4NBLp:	mov	cx, 10000	
Wait4NotBusy:	in test	al, dx al, 80h	;Read status register value. ;Bit 7 = 1 if busy.
	loopne	Wait4NotBusy	Repeat while busy, 10000 times.
	je	ItsNotbusy	;Leave loop if not busy.
	call jmp	TestAbort W4NBLp	;Check for Ctrl-C.
; Okay, put th	he data on	the data lines:	
ItsNotBusy:	dec	dx	;Point at data register.
reshoebusy.	mov	al, ah	;Get a copy of the data.
	and	al, OFh	;Strip out H.O. nibble
	out or	dx, al al, 10h	;"Prime" data lines, data not avail. ;Turn data available on.
	out	dx, al	;Send data w/data available strobe.
	ctrl-C so		eiving site. Every now and then wort the transmission program from
	inc	dx	;Point at status register.
W4ALp:	mov	cx, 10000	;Times to loop between ctrl-C checks.
Wait4Ack:	in test	al, dx al, 80h	;Read status port. ;Ack = 1 when rcvr acknowledges.
	loope	Wait4Ack	;Repeat 10000 times or until ack.
	jne	GotAck TostAbort	;Branch if we got an ack.
	call	TestAbort	;Every 10000 calls, check for a

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imp	W4ALp	; ctrl-C from the user.
mp	W4ALD	; CLII-C IIOM LNE USEL.

; Send the data not available signal to the receiver:

GotAck:	dec	dx	;Point at data register.
	mov	al, 0	;Write a zero to bit 4, this appears
	out	dx, al	; as a one in the rcvr's busy bit.

; Okay, on to the H.O. nibble:

s.

; Okay, put the data on the data lines:

NotBusy2:	dec mov shr out or	dx al, ah al, 4 dx, al al, 10h dx, al	;Point at data register. ;Retrieve data to get H.O. nibble. ;Move H.O. nibble to L.O. nibble. ;"Prime" data lines. ;Data + data available strobe. :Send data w/data available strobe
	out	dx, al	;Send data w/data available strobe.

; Wait for the acknowledge from the receiving site:

W4A2Lp:	inc mov	dx cx, 10000	;Point at status register.
Wait4Ack2:	in test loope jne call jmp	al, dx al, 80h Wait4Ack2 GotAck2 TestAbort W4A2Lp	<pre>;Read status port. ;Ack = 1 ;While while no acknowledge ;H.O. bit = 1 (ack)? ;Check for ctrl-C</pre>

; Send the data not available signal to the receiver:

GotAck2:	dec	dx	;Point at data register.
	mov	al, O	;Output a zero to bit #4 (that
	out	dx, al	; becomes busy=1 at rcvr).
SendByte	mov pop pop ret endp	al, ah dx cx	;Restore original data in AL.

; Synchroniza ;			
; Send0s-	Transmits a zero to the receiver site and then waits to		
;	see if it gets a set of ones back. Returns carry set if		
;	this works, returns carry clear if we do not get a set of		
;	ones back in a reasonable amount of time.		
Send0s	proc push push	near cx dx	
	mov	dx, MyPortAdrs	
	mov	al, O	;Write the initial zero
	out	dx, al	; value to our output port.
Wait41s:	xor	cx, cx	;Checks for ones 10000 times.
	inc	dx	;Point at status port.
	in	al, dx	;Read status port.
	dec	dx	;Point back at data port.

	and cmp loopne je clc pop pop ret	al, 78h al, 78h Wait41s Got1s dx cx	;Mask input bits. ;All ones yet? ;Branch if success. ;Return failure.
Gotls:	stc pop pop ret	dx cx	;Return success.
Send0s	endp		
; Sendls- ; ;	see if it this works	gets a set of zero	eceiver site and then waits to os back. Returns carry set if lear if we do not get a set of amount of time.
Send1s	proc push push	near cx dx	
	mov	dx, MyPortAdrs	;LPT1: base address.
	mov out	al, OFh dx, al	;Write the "all ones" ; value to our output port.
Wait40s:	mov inc dec and loopne je clc pop pop ret	cx, 0 dx al, dx dx al, 78h Wait40s Got0s dx cx	<pre>;Point at input port. ;Read the status port. ;Point back at data port. ;Mask input bits. ;Loop until we get zero back. ;All zeros? If so, branch. ;Return failure.</pre>
Got0s:	stc pop pop ret	dx cx	;Return success.
Sendls	endp		

; Synchronize- This procedure slowly writes all zeros and all ones to its ; output port and checks the input status port to see if the ; receiver site has synchronized. When the receiver site ; is synchronized, it will write the value 05h to its output ; port. So when this site sees the value 05h on its input ; port, both sites are synchronized. Returns with the ; carry flag set if this operation is successful, clear if ; unsuccessful.

Synchronize	proc print	near	
	byte byte	"Synchronizing wi cr,lf,0	th receiver program"
	mov	dx, MyPortAdrs	
SyncLoop:	mov call jc	cx, TimeOutConst SendOs Gotls	;Time out delay. ;Send zero bits, wait for ; ones (carry set=got ones).

; If we didn't get what we wanted, write some ones at this point and see ; if we're out of phase with the receiving site.

Retrv0: call Send1s ;Send ones, wait for zeros. SvncLoop ;Carry set = got zeros. ic ; Well, we didn't get any response yet, see if the user has pressed ctrl-C ; to abort this program. DoRetrv: call TestAbort : Okay, the receiving site has yet to respond. Go back and try this again. loop SyncLoop ; If we've timed out, print an error message and return with the carry ; flag clear (to denote a timeout error). print "Transmit: Timeout error waiting for receiver" byte cr, lf, 0 bvte clc ret : Okay, we wrote some zeros and we got some ones. Let's write some ones ; and see if we get some zeros. If not, retry the loop. Got1s: Send1s call ;Send one bits, wait for inc DoRetrv ; zeros (carry set=got zeros). ; Well, we seem to be synchronized. Just to be sure, let's play this out ; one more time. call Send0s ;Send zeros, wait for ones. Retry0 jnc call Send1s ;Send ones, wait for zeros. jnc DoRetry ; We're syncronized. Let's send out the 05h value to the receiving ; site to let it know everything is cool: al, 05h ;Send signal to receiver to mov ; tell it we're sync'd. out dx, al ;Long delay to give the rcvr xor CX, CX FinalDelav: loop FinalDelay ; time to prepare. print byte "Synchronized with receiving site" byte cr, lf, 0 stc ret. Synchronize endp ; File I/O routines: GetFileInfo-Opens the user specified file and passes along the file ; name and file size to the receiving site. Returns the ; carry flag set if this operation is successful, clear if ; unsuccessful. ; GetFileInfo proc near ; Get the filename from the DOS command line: mov ax, 1 arqv word ptr FileNamePtr, di mov mov word ptr FileNamePtr+2, es printf "Opening %^s\n",0 byte dword FileNamePtr

; Open the file:

push ds mov ax, 3D00h ;Open for reading. lds dx, FileNamePtr int 21h pop ds jc BadFile mov FileHandle, ax

; Compute the size of the file (do this by seeking to the last position ; in the file and using the return position as the file length):

mov	bx, ax	;Need handle in BX.
mov	ax, 4202h	;Seek to end of file.
xor	CX, CX	;Seek to position zero
xor	dx, dx	; from the end of file.
int	21h	
jc	BadFile	

; Save final position as file length:

mov	word ptr FileSize, ax	
mov	word ptr FileSize+2,	dx

; Need to rewind file back to the beginning (seek to position zero):

mov	bx, FileHandle	;Need handle in BX.
mov	ax, 4200h	;Seek to beginning of file.
xor	CX, CX	;Seek to position zero
xor	dx, dx	
int	21h	
jc	BadFile	

; Okay, transmit the good stuff over to the receiving site:

	mov call mov call mov call mov call	al, byte ptr FileSize SendByte al, byte ptr FileSize+1 SendByte al, byte ptr FileSize+2 SendByte al, byte ptr FileSize+3 SendByte	;Send the file ; size over.
SendName:	les mov call inc cmp jne stc ret	bx, FileNamePtr al, es:[bx] SendByte bx al, 0 SendName	;Send the characters ; in the filename to ; the receiver until ; we hit a zero byte. ;Return success.
BadFile: GetFileInfo	print byte puti putcr clc ret endp	<pre>"Error transmitting file inform</pre>	nation:",0

; GetFileData-This procedure reads the data from the file and transmits ; it to the receiver a byte at a time.

GetFileData	proc	near	
	mov	ah, 3Fh	;DOS read opcode.
	mov	cx, 512	;Read 512 bytes at a time.
	mov	bx, FileHandle	;File to read from.
	lea	dx, FileBuffer	;Buffer to hold data.
	int	21h	;Read the data

	jc	GFDError	;Quit if error reading data.	
XmitLoop:	mov jcxz lea mov call inc	cx, ax GFDDone bx, FileBuffer al, [bx] SendByte bx	;Save # of bytes actually read. ; quit if at EOF. ;Send the bytes in the file ; buffer over to the rcvr ; one at a time.	
	loop jmp	XmitLoop GetFileData	;Read rest of file.	
GFDError:	print byte puti print	"DOS error #",	0	
GFDDone: GetFileData	byte ret endp	" while readin	g file",cr,lf,0	
; Okay, here'	s the main	program that co	ntrols everything.	
Main	proc mov mov meminit	ax, dseg ds, ax		
; First, get	the address	of LPT1: from	the BIOS variables area.	
	mov mov mov mov	ax, 40h es, ax ax, es:[PrtrBa MyPortAdrs, ax		
; See if we h	ave a filer	ame parameter:		
	argc cmp je print byte jmp	cx, 1 GotName "Usage: transm Quit	it <filename>",cr,lf,0</filename>	
GotName:	call jnc	Synchronize Quit	;Wait for the transmitter program.	
	call jnc	GetFileInfo Quit	;Get file name and size.	
	call	GetFileData	;Get the file's data.	
Quit: Main	ExitPgm endp		;DOS macro to quit program.	
cseg	ends			
sseg stk sseg	segment byte ends	para stack `st 1024 dup ("sta		
zzzzzzseg LastBytes zzzzzseg	segment byte ends end	para public `zzzzzz' 16 dup (?) Main		

Here is the receiver program that accepts and stores away the data sent by the program above:

; RECEIVE.ASM ; This program is the receiver portion of the programs that transmit files ; across a Laplink compatible parallel cable. ; This program assumes that the user want to use LPT1: for transmission. Adjust the equates, or read the port from the command line if this ; is inappropriate. .286 .xlist include stdlib.a includelib stdlib.lib .list seament para public 'data' dsea TimeOutConst equ 100 ;About 1 min on 66Mhz 486. PrtrBase 8 ;Offset to LPT1: adrs. equ MvPortAdrs ? word ;Holds printer port address. FileHandle word 2 ;Handle for output file. FileBuffer byte 512 dup (?) ;Buffer for incoming data. FileSize dword ;Size of incoming file. ? FileName 128 dup (0) ;Holds filename bvte ends dsea segment para public 'code' csea assume cs:cseq, ds:dseq ; TestAbort-Reads the keyboard and gives the user the opportunity to hit the ctrl-C key. ; TestAbort proc near push ax mov ah, 1 16h int ;See if keypress. NoKeypress ie ;Read char, chk for ctrl-C ah, 8 mov int 21h NoKeyPress: ax pop ret TestAbort endp ; GetByte-Reads a single byte from the parallel port (four bits at at time). Returns the byte in AL. ; GetByte proc near push СХ push dx ; Receive the L.O. Nibble. mov dx, MyPortAdrs al, 10h ;Signal not busy. mov out dx, al ; Point at status port inc dx W4DLp: cx, 10000 mov Wait4Data: in al, dx ;See if data available. al, 80h ; (bit 7=0 if data available). test Wait4Data loopne ;Is data available? DataIsAvail je call TestAbort ; If not, check for ctrl-C.

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	jmp	W4DLp	
DataIsAvail:	shr and mov	al, 3 al, OFh ah, al	;Save this four bit package ; (This is the L.O. nibble ; for our byte).
	dec mov out	dx al, 0 dx, al	;Point at data register. ;Signal data taken.
W4ALp: Wait4Ack:	inc mov in test loope jne call jmp	dx cx, 10000 al, dx al, 80h Wait4Ack NextNibble TestAbort W4ALp	;Point at status register. ;Wait for transmitter to ; retract data available. ;Loop until data not avail. ;Branch if data not avail. ;Let user hit ctrl-C.
; Receive the	H.O. nibbl	e:	
NextNibble:	dec mov out inc	dx al, 10h dx, al dx	;Point at data register. ;Signal not busy ;Point at status port
W4D2Lp: Wait4Data2:	mov in test loopne je call jmp	cx, 10000 al, dx al, 80h Wait4Data2 DataAvail2 TestAbort W4D2Lp	;See if data available. ; (bit 7=0 if data available). ;Loop until data available. ;Branch if data available. ;Check for ctrl-C.
DataAvail2:	shl and or dec mov out	al, 1 al, 0F0h ah, al dx al, 0 dx, al	;Merge this H.O. nibble ; with the existing L.O. ; nibble. ;Point at data register. ;Signal data taken.
W4A2Lp: Wait4Ack2:	inc mov in test loope jne call jmp	dx cx, 10000 al, dx al, 80h Wait4Ack2 ReturnData TestAbort W4A2Lp	;Point at status register. ;Wait for transmitter to ; retract data available. ;Wait for data not available. ;Branch if ack. ;Check for ctrl-C
ReturnData:	mov pop pop	al, ah dx cx	;Put data in al.
GetByte	ret endp		

; Synchronize- ; ; ; ;	This procedure waits until it sees all zeros on the input bits we receive from the transmitting site. Once it receives all zeros, it writes all ones to the output port. When all ones come back, it writes all zeros. It repeats this process until the transmitting site writes the value 05h.			
Synchronize	proc print byte byte	<pre>near "Synchronizing with transmitter program" cr,lf,0</pre>		
	mov mov out mov	dx, MyPortAdrs al, 0 dx, al bx, TimeOutConst	;Initialize our output port ; to prevent confusion. ;Time out condition.	

SyncLoop: SyncLoop0:	mov inc in dec and cmp je cmp loopne	cx, 0 dx al, dx dx al, 78h al, 78h Got1s al, 0 SyncLoop0	<pre>;For time out purposes. ;Point at input port. ;Read our input bits. ;Keep only the data bits. ;Check for all ones. ;Branch if all ones. ;See if all zeros.</pre>	
; Since we ju	st saw a ze	ero, write all ones	to the output port.	
	mov out	al, OFFh ;W dx, al	rite all ones	
; Now wait fo	r all ones	to arrive from the	e transmitting site.	
SyncLoop1:	inc in dec and cmp loopne je	dx al, dx dx al, 78h al, 78h SyncLoopl Gotls	;Point at status register. ;Read status port. ;Point back at data register. ;Keep only the data bits. ;Are they all ones? ;Repeat while not ones. ;Branch if got ones.	
; If we've ti ; abort.	med out, ch	eck to see if the	user has pressed ctrl-C to	
	call dec jne	TestAbort bx SyncLoop	;Check for ctrl-C. ;See if we've timed out. ;Repeat if time-out.	
	print byte byte clc ret	"Receive: connect cr,lf,0	tion timed out during synchronization";Signal time-out.	
	ions until		ero and a one. Send the two the transmitting site or the	
Got1s:	inc in dec shr and cmp je not out call jmp	dx al, dx dx al, 3 al, OFh al, OFh al, O5h Synchronized al dx, al TestAbort Got1s	<pre>;Point at status register. ;Just copy whatever appears ; in our input port to the ; output port until the ; transmitting site sends ; us the value 05h ;Keep inverting what we get ; and send it to xmitter. ;Check for CTRL-C here.</pre>	
; Okay, we're synchronized. Return to the caller.				
Synchronized: Synchronize	and out print byte byte stc ret endp	al, OFh dx, al "Synchronized wit cr,lf,0	;Make sure busy bit is one ; (bit 4=0 for busy=1). th transmitting site"	
-	- The transm		nds us the file length and a	
;	zero term	inated filename. Ge	et that data here.	
GetFileInfo	proc mov mov	near dx, MyPortAdrs al, 10h	;Set busy bit to zero.	

out dx, al :Tell xmit pom, we're ready. ; First four bytes contain the filesize: call GetByte byte ptr FileSize, al mov GetByte call byte ptr FileSize+1, al mov call GetByte mov byte ptr FileSize+2, al call GetBvte byte ptr FileSize+3, al mov ; The next n bytes (up to a zero terminating byte) contain the filename: bx, 0 mov GetFileName: GetByte call FileName[bx], al mov TestAbort call inc bx cmp al, 0 jne GetFileName ret GetFileInfo endp ; GetFileData-Receives the file data from the transmitting site and writes it to the output file. ; GetFileData proc near ; First, see if we have more than 512 bytes left to go word ptr FileSize+2, 0 amp ; If H.O. word is not MoreThan512 ; zero, more than 512. ine word ptr FileSize, 512 cmp ; If H.O. is zero, just jbe LastBlock ; check L.O. word. ; We've got more than 512 bytes left to go in this file, read 512 bytes ; at this point. cx, 512 bx, FileBuffer MoreThan512: mov ;Receive 512 bytes ; from the xmitter. lea GetBvte ;Read a byte. ReadLoop: call mov [bx], al ; Save the byte away. inc bx ;Move on to next loop ReadLoop ; buffer element. ; Okay, write the data to the file: ah, 40h mov ;DOS write opcode. mov bx, FileHandle ;Write to this file. cx, 512 dx, Filebuffer ;Write 512 bytes. mov ;From this address. lea int 21h BadWrite ;Quit if error. jс ; Decrement the file size by 512 bytes: word ptr FileSize, 512 sub ;32-bit subtraction word ptr FileSize, 0 sbb ; of 512. jmp GetFileData ; Process the last block, that contains 1..511 bytes, here. LastBlock: mov cx, word ptr FileSize ;Receive the last ; 1..511 bytes from lea bx, FileBuffer ReadLB: call GetByte ; the transmitter. [bx], al mov inc bx ReadLB loop

	mov mov lea int jnc	ah, 40h bx, FileHandle cx, word ptr File dx, Filebuffer 21h Closefile	eSize	;Write the last block ; of bytes to the ; file.
BadWrite:	print byte puti print byte	"DOS error #",0	data.",cr,lf,0	
; Close the t	file here.			
CloseFile:	mov mov int ret	bx, FileHandle ah, 3Eh 21h		;Close this file. ;DOS close opcode.
GetFileData	endp			
	main progra	am that gets the wh	ole ball rolli	ng.
Main	proc mov mov meminit	ax, dseg ds, ax		
; First, get	the address	s of LPT1: from the	e BIOS variable	s area.
	mov mov mov mov	ax, 40h es, ax ax, es:[PrtrBase] MyPortAdrs, ax		98 variable segment.
	call jnc	Synchronize Quit	;Wait for the	e transmitter program.
	call	GetFileInfo	;Get file nam	e and size.
	printf byte dword	"Filename: %s\nFi Filename, FileSiz	%s\nFile size: %ld\n″,0 ileSize	
	mov mov lea int jnc	ah, 3Ch cx, 0 dx, Filename 21h GoodOpen	;Create file. ;Standard att	
	print byte jmp	"Error opening fi Quit	ile",cr,lf,0	
GoodOpen:	mov call	FileHandle, ax GetFileData	;Get the file	's data.
Quit: Main	ExitPgm endp		;DOS macro to	quit program.
cseg	ends			
sseg stk sseg	segment byte ends	para stack `stac} 1024 dup ("stack		
zzzzzseg LastBytes zzzzzseg	segment byte ends	para public `zzzzz' 16 dup (?)		
ر	end	Main		

## 21.5 Summary

The PC's parallel port, though originally designed for controlling parallel printers, is a general purpose eight bit output port with several handshaking lines you can use to control many other devices in addition to printers.

In theory, parallel communications should be many times faster than serial communications. In practice, however, real world constraints and economics prevent this from being the case. Nevertheless, you can still connect high performance devices to the PC's parallel port.

The PC's parallel ports come in two varieties: unidirectional and bidirectional. The bidirectional versions are available only on PS/2s, certain laptops, and a few other machines. Whereas the eight data lines are output only on the unidirectional ports, you can program them as inputs or outputs on the bidirectional port. While this bidirectional operation is of little value to a printer, it can improve the performance of other devices that connect to the parallel port, such as disk and tape drives, network adapters, SCSI adapters, and so on.

When the system communicates with some other device over the parallel port, it needs some way to tell that device that data is available on the data lines. Likewise, the devices needs some way to tell the system that it is not busy and it has accepted the data. This requires some additional signals on the parallel port known as handshaking lines. A typical PC parallel port provides three handshaking signals: the data available strobe, the data taken acknowledge signal, and the device busy line. These lines easily control the flow of data between the PC and some external device.

In addition to the handshaking lines, the PC's parallel port provides several other auxiliary I/O lines as well. In total, there are 12 output lines and five input lines on the PC's parallel port. There are three I/O ports in the PC's address space associated with each I/O port. The first of these (at the port's base address) is the data register. This is an eight bit output register on unidirectional ports, it is an input/output register on bidirectional ports. The second register, at the base address plus one, is the status register. The status register is an input port. Five of those bits correspond to the five input lines on the PC's parallel port. The third register (at base address plus two) is the control register. Four of these bits correspond to the additional four output bits on the PC, one of the bits controls the IRQ line on the parallel port, and a sixth bit controls the data direction on the birdirectional ports.

For more information on the parallel port's hardware configuration, see:

- "Basic Parallel Port Information" on page 1199
- "The Parallel Port Hardware" on page 1201

Although many vendors use the parallel port to control lots of different devices, a parallel printer is still the device most often connected to the parallel port. There are three ways application programs commonly send data to the printer: by calling DOS to print a character, by calling BIOS' int 17h ISR to print a character, or by talking directly to the parallel port. You should avoid this last technique because of possible software incompatibilities with other devices that connect to the parallel port. For more information on printing data, including how to write your own int 17h ISR/printer driver, see:

- "Controlling a Printer Through the Parallel Port" on page 1202
- "Printing via DOS" on page 1203
- "Printing via BIOS" on page 1203
- "An INT 17h Interrupt Service Routine" on page 1203

One popular use of the parallel port is to transfer data between two computers; for example, transferring data between a desktop and a laptop machine. To demonstrate how to use the parallel port to control other devices besides printers, this chapter presents a program to transfer data between computers on the unidirectional parallel ports (it also works on bidirectional ports). For all the details, see

• "Inter-Computer Communications on the Parallel Port" on page 1209